ProGraML: Graph-based Deep Learning for **Program Optimization** and Analysis. **Chris Cummins Facebook AI Research**

"machine learning for compilers for machine learning"

Compilers

Machine Learning

Tuning optimizing compilers...

The problem

- 1000s of variables
- Limited by domain expertise
- Compiler / HW keeps changing

The cost

- Bad heuristics
- Wasted energy, \$\$\$
- Widening performance gap

"Build an optimizing compiler, your code will be fast for a day. Teach a compiler to optimize ... "



Summarize the program

Program



Collect examples

Features



Learn from examples

Best Param

...









Why aren't our compilers full of



Learning without features (Cummins et al., PACT 17) "End-to-end Deep Learning of **Optimization Heuristics**" 1. Input kernel void A(global float* a, const float b) { a[get_global_id(0)] *= 3.14 + b; 2. Vocab Token Index Token Index kernel 0 9 , [space] 10 const void 2 11 b 3 12 Α 181 tokens 13 4 global 5 14 \n float 6 15 Optimization 7 16 * get_global_id Decision 17 8 0 а 3. Encoded 0 3 5

LSTM

The problem with code representations

Source code is *highly structured*

It isn't a vector of numbers

Feature vectors are easy to fool (e.g. insert dead code).



It isn't a sequence of tokens

Sequential representations fail on non-linear relations, long-range deps.



Can we make ML think like a compiler?

Program Graphs for Machine Learning

General-purpose representation of programs for optimization tasks.

Task independent - capture structured relations fundamental to program reasoning (i.e. data flow analysis)

Language independent - derived from compiler IRs

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	ABSTRACT	

Building ProGraML: IR

Derive IR from input program (here, LLVM)

Why IR?

Language **agnostic**

(e.g. C, C++, OpenCL, Swift, Haskell, Java for LLVM)

We want to improve compiler decisions, so use a **compiler's eye** view.





Building ProGraML: Control-flow

Full-flow-graph: represent each instruction as a vertex.

Vertex label is the instruction name.

Edges are **control-flow**.

Edge position attribute for **branching control-flow**.



Building ProGraML: Data-flow

Add graph vertices for constants (diamonds) and variables (oblongs).

Edges are data-flow.

Edge position attribute for **operand order**.



Building ProGraML: Call-flow

Edges are **call-flow**.

Inbound edge to function entry instruction.

Outbound edge from (all) **function exit** instruction(s).



Building ProGraML: Types

Nodes represent **types**, Edges are **instances**.

Types are **composable**. Edge position per field.





Learning with ProGraML: Node Embeddings

Use vertex labels as embedding keys $br \rightarrow 0$ add $\rightarrow 1$ (i32) $\rightarrow 2$

Derive vocab from set of unique vertex labels on training graphs.

Separate type/instruction nodes leads to **compact vocab**, excellent coverage on unseen programs compared to prior approaches:

	Vocabulary size	Test coverage
inst2vec [12]	8,565	34.0%
CDFG [14]	75	47.5%
ProGRAML	2,230	98.3% *without types

inst2vec: combined instruction+operands <u>CDFG</u>: uses only instructions for vocab, ignores data

Learning with ProGraML: GGNNs

Message Passing

$$M(h_w^{t-1}, e_{wv}) = W_{\text{type}(e_{wv})} \left(h_w^{t-1} \odot p(e_{wv}) \right) + b_{\text{type}(e_{wv})}$$

6 typed weight matrices for {forwards,backwards} {control,data,call} edge types

Position gating to differentiate control branches and operand order

Readout Head

$$R_{v}(h_{v}^{T},h_{v}^{0}) = \sigma\left(f(h_{v}^{T},h_{v}^{0})\right) \cdot g(h_{v}^{T}$$
 prediction after **T**

per-vertex prediction after message-passing steps

Deep Data Flow

Dataset: 450k LLVM-IRs covering 5 programming languages

F1 scores

CDEC

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inct?voo

	11312460	ODFG	FIGUAINL
Reachability Trivial forwards control-flow E.g. dead code elimination	0.012	0.998	0.998
Dominance Forwards control-flow E.g. global code motion	0.004	0.999	1.000
Data Dependencies Forwards data-flow E.g. instruction selection	-	-	0.997
Live-out Variables Backwards control- and data-flow E.g. register allocation	-	-	0.937
Global Common Subexpressions Instruction/operand sensitive E.g. GCS Elimination	0.000	0.009	0.996

Deep Data Flow

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Caveat: limited problem size

Data flow analyses iterate until a fixed point is reached.

GGNNs iterate for a fixed number of timesteps **T**.

For each example in the train/test sets, we count the number of steps required for an iterative analysis to solve.

We then filter the train/test set to include only examples which the iterative analysis required **<= T** steps to solve.

Previous slide was **T=30**, excluding 28.7% of examples.

Next slide shows performance models, trained on **T=30**, with different inference steps (**T=60**, **T=200**).



Scaling to larger problems

F1 scores

30 200 Dataset: 450k LLVM-IRs covering 5 programming languages timesteps 60 timesteps timesteps Reachability 0.998 0.997 0.943 Trivial forwards control-flow E.g. dead code elimination Dominance 1.000 0.9910.123 Forwards control-flow E.g. global code motion **Data Dependencies** 0.997 0.993 0.965 Forwards data-flow E.g. instruction selection Live-out Variables 0.937 0.625 0.939 Backwards control- and data-flow E.g. register allocation **Global Common Subexpressions** 0.996 0.967 0.959 Instruction/operand sensitive E.g. GCS Elimination



Downstream tasks





1.35× improvement over state-of-art

2. Heterogeneous Device Mapping



1.20× improvement over state-of-art

Further Reading



https://arxiv.org/abs/2003.10536

In-browser demo https://chriscummins.cc/s/program_explorer Source code + datasets https://github.com/ChrisCummins/ProGraML Apache 2.0

Conclusions

Reasoning about programs requires the right combination of representation + model.

ProGraML: combines control-, data-, call-, and type-graphs to model programs at IR level.

When processed with GGNNs, significantly outperforms prior approaches.

Interesting challenges

1. Processing arbitrary sized graphs.

Idea: Structure the MPNN like an iterative DF solver, self-terminating.

- 2. Handling **unbounded vocabularies**, e.g. compound types or MLIR dialects. Idea: decompose types into tree structure in graph.
- 3. Representing literal values.

Requires new vocabulary encoding.